

# Convergence Rates in the Probabilistic Analysis of Algorithms

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## Abstract

In this extended abstract a general framework is developed to bound rates of convergence for sequences of random variables as they mainly arise in the analysis of random trees and divide-and-conquer algorithms. The rates of convergence are bounded in the Zolotarev distances. Concrete examples from the analysis of algorithms and data structures are discussed as well as a few examples from other areas. They lead to convergence rates of polynomial and logarithmic order. Our results show how to obtain a significantly better bound for the rate of convergence when the limiting distribution is Gaussian.

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## 1 Introduction and notation

In this extended abstract we consider a general recurrence for (probability) distributions which covers many instances of complexity measures of divide-and-conquer algorithms and parameters of random search trees. We consider a sequence  $(Y_n)_{n \geq 0}$  of  $d$ -dimensional random vectors satisfying the distributional recursion

$$Y_n \stackrel{d}{=} \sum_{r=1}^K A_r(n) Y_{I_r^{(n)}}^{(r)} + b_n, \quad n \geq n_0, \quad (1)$$

where  $(A_1(n), \dots, A_K(n), b_n, I^{(n)})$ ,  $(Y_n^{(1)})_{n \geq 0}, \dots, (Y_n^{(K)})_{n \geq 0}$  are independent, the coefficients  $A_1(n), \dots, A_K(n)$  are random  $(d \times d)$ -matrices,  $b_n$  is a  $d$ -dimensional random vector,  $I^{(n)} = (I_1^{(n)}, \dots, I_K^{(n)})$  is a random vector in  $\{0, \dots, n\}^K$ ,  $n_0 \geq 1$  and  $(Y_n^{(r)})_{n \geq 0} \stackrel{d}{=} (Y_n)_{n \geq 0}$  for  $r = 1, \dots, K$ . Moreover,  $K \geq 1$  is a fixed integer, but extensions to  $K$  being random and depending on  $n$  are possible.

This is the framework of [14] where some general convergence results are shown for appropriate normalizations of the  $Y_n$ . The content of the present extended abstract is to also study the rates of convergence in such limit theorems.

We define the normalized sequence  $(X_n)_{n \geq 0}$  by

$$X_n := C_n^{-1/2}(Y_n - M_n), \quad n \geq 0,$$

where  $M_n$  is a  $d$ -dimensional vector and  $C_n$  a positive definite  $(d \times d)$ -matrix. Essentially, we choose  $M_n$  as the mean and  $C_n$  as the covariance matrix of  $Y_n$  if they exist or as the leading order terms in expansions of these moments as  $n \rightarrow \infty$ . The normalized quantities satisfy the following modified recursion:

$$X_n \stackrel{d}{=} \sum_{r=1}^K A_r^{(n)} X_{I_r^{(n)}}^{(r)} + b^{(n)}, \quad n \geq n_0, \quad (2)$$

with

$$A_r^{(n)} := C_n^{-1/2} A_r(n) C_{I_r^{(n)}}^{1/2}, \quad b^{(n)} := C_n^{-1/2} \left( b_n - M_n + \sum_{r=1}^K A_r(n) M_{I_r^{(n)}} \right) \quad (3)$$

and independence relations as in (1).

In the context of the contraction method the aim is to establish transfer theorems of the following form: After verifying the assumptions of appropriate convergence of the coefficients  $A_r^{(n)} \rightarrow A_r^*$ ,  $b^{(n)} \rightarrow b^*$  then convergence in distribution of random vectors  $(X_n)$  to a limit  $X$  is implied. The limit distribution  $\mathcal{L}(X)$  is identified by a fixed-point equation obtained from (2) by considering formally  $n \rightarrow \infty$ :

$$X \stackrel{d}{=} \sum_{r=1}^K A_r^* X^{(r)} + b^*.$$

Here  $(A_1^*, \dots, A_K^*, b^*)$ ,  $X^{(1)}, \dots, X^{(K)}$  are independent and  $X^{(r)} \stackrel{d}{=} X$  for  $r = 1, \dots, K$ .

The aim of the present extended abstract is to endow such general transfer theorems with bounds on the rates of convergence. As a distance measure between (probability) distributions we use the Zolotarev metric. For various of the applications we discuss, bounds on the rate of convergence have been derived one by one for more popular distance measures such as the Kolmogorov–Smirnov distance. However, the transfer theorems of the present paper in terms of the smoother Zolotarev metrics are easy to apply and cover a broad range of applications at once. A crucial role is played by a factor 3 in the exponent of these orders in cases where the normal distribution is the limiting distribution, see Remark 4.

In the rest of this section we fix some notation. Regarding norms of vectors and (random) matrices we denote for  $x \in \mathbb{R}^d$  by  $\|x\|$  its Euclidean norm and for a random vector  $X$  and some  $0 < p < \infty$ , we set  $\|X\|_p := \mathbb{E}[\|X\|^p]^{(1/p) \wedge 1}$ . Furthermore, for a  $(d \times d)$ -matrix  $A$ ,  $\|A\|_{\text{op}} := \sup_{\|x\|=1} \|Ax\|$  denotes the spectral norm of  $A$  and for a random such  $A$  we define  $\|A\|_p := \mathbb{E}[\|A\|_{\text{op}}^p]^{(1/p) \wedge 1}$  for a random square matrix and  $0 < p < \infty$ . Note that for a symmetric  $(d \times d)$ -matrix  $A$ , we have  $\|A\|_{\text{op}} = \max\{|\lambda| : \lambda \text{ eigenvalue of } A\}$ . By  $\text{Id}_d$  the  $d$ -dimensional unit matrix is denoted. For multilinear forms the norm is defined similarly.

Furthermore we define by  $\mathcal{P}^d$  the space of probability distributions in  $\mathbb{R}^d$  (endowed with the Borel  $\sigma$ -field), by  $\mathcal{P}_s^d := \{\mathcal{L}(X) \in \mathcal{P}^d : \|X\|_s < \infty\}$  and for a vector  $m \in \mathbb{R}^d$ , and a symmetric positive semidefinite  $(d \times d)$ -matrix  $C$  the spaces

$$\begin{aligned} \mathcal{P}_s^d(m) &:= \{\mathcal{L}(X) \in \mathcal{P}_s^d : \mathbb{E}[X] = m\}, \quad s > 1, \\ \mathcal{P}_s^d(m, C) &:= \{\mathcal{L}(X) \in \mathcal{P}_s^d : \mathbb{E}[X] = m, \text{Cov}(X) = C\}, \quad s > 2. \end{aligned} \quad (4)$$

We use the convention  $\mathcal{P}_s^d(m) := \mathcal{P}_s^d$  for  $s \leq 1$  and  $\mathcal{P}_s^d(m, C) := \mathcal{P}_s^d(m)$  for  $s \leq 2$ .

The Zolotarev metrics  $\zeta_s$ , [19], are defined for probability distributions  $\mathcal{L}(X), \mathcal{L}(Y) \in \mathcal{P}^d$  by

$$\zeta_s(X, Y) := \zeta_s(\mathcal{L}(X), \mathcal{L}(Y)) = \sup_{f \in \mathcal{F}_s} |E(f(X) - f(Y))|,$$

where for  $s = m + \alpha$ ,  $0 < \alpha \leq 1$ ,  $m \in \mathbb{N}_0$ ,

$$\mathcal{F}_s := \{f \in C^m(\mathbb{R}^d, \mathbb{R}) : \|f^{(m)}(x) - f^{(m)}(y)\| \leq \|x - y\|^\alpha\}.$$

Note that these distance measures may be infinite. Finite metrics are given by  $\zeta_s$  on  $\mathcal{P}_s^d$  for  $0 \leq s \leq 1$ , by  $\zeta_s$  on  $\mathcal{P}_s^d(m)$  for  $1 < s \leq 2$ , and by  $\zeta_s$  on  $\mathcal{P}_s^d(m, C)$  for  $2 < s \leq 3$ , cf. (4).

## 2 Results

We return to the situation outlined in the introduction, where we have normalized  $(Y_n)_{n \geq 0}$  in the following way:

$$X_n := C_n^{-1/2}(Y_n - M_n), \quad n \geq 0, \quad (5)$$

where  $M_n$  is a  $d$ -dimensional random vector and  $C_n$  a positive definite  $(d \times d)$ -matrix. As recalled in Section 1, for  $s > 1$ , we may fix the mean and covariance matrix of the scaled quantities to guarantee the finiteness of the  $\zeta_s$ -metric. Therefore, we choose  $M_n = \mathbb{E}[Y_n]$  for  $n \geq 0$  and  $s > 1$ . For  $s > 2$ , we additionally have to control the covariances of  $X_n$ . We assume that there exists an  $n_1 \geq 0$  such that  $\text{Cov}(Y_n)$  is positive definite for  $n \geq n_1$  and choose  $C_n = \text{Cov}(Y_n)$  for  $n \geq n_1$  and  $C_n = \text{Id}_d$  for  $n < n_1$ . For  $s \leq 2$ , we just assume that  $C_n$  is positive definite and set  $n_1 = 0$  in this case.

The normalized quantities satisfy the modified recursion

$$X_n \stackrel{d}{=} \sum_{r=1}^K A_r^{(n)} X_{I_r^{(n)}}^{(r)} + b^{(n)}, \quad n \geq n_0,$$

with  $A_r^{(n)}$  and  $b^{(n)}$  given in (3). The following theorem discusses a general framework to bound rates of convergence for the sequence  $(X_n)_{n \geq 0}$ . For the proof, we need some technical conditions which guarantee that the sizes  $I_r^{(n)}$  of the subproblems grow with  $n$ . More precisely, we will assume that there exists some monotonically decreasing sequence  $R(n) > 0$  with  $R(n) \rightarrow 0$  such that

$$\|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_s = O(R(n)), \quad n \rightarrow \infty, \quad (6)$$

for all  $\ell \in \mathbb{N}$  and  $r = 1, \dots, K$  and that

$$\|\mathbf{1}_{\{I_r^{(n)} = n\}} A_r^{(n)}\|_s \rightarrow 0, \quad n \rightarrow \infty, \quad (7)$$

for all  $r = 1, \dots, K$ .

### 2.1 A general transfer theorem for rates of convergence

Our first result is a direct extension of the main Theorem 4.1 in [14], where we essentially only make all the estimates there explicit. The main result of the present extended abstract is contained in Section 2.2.

► **Theorem 1.** *Let  $(X_n)_{n \geq 0}$  be  $L_s$ -integrable,  $0 < s \leq 3$ , and satisfy recurrence (5) with the choices for  $M_n$  and  $C_n$  specified there. We assume that there exist  $s$ -integrable  $A_1^*, \dots, A_K^*, b^*$  and some monotonically decreasing sequence  $R(n) > 0$  with  $R(n) \rightarrow 0$  such that, as  $n \rightarrow \infty$ ,*

$$\|b^{(n)} - b^*\|_s + \sum_{r=1}^K \|A_r^{(n)} - A_r^*\|_s = O(R(n)). \quad (8)$$

If conditions (6) and (7) are satisfied and if

$$\limsup_{n \rightarrow \infty} \mathbb{E} \sum_{r=1}^K \left( \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^s \right) < 1, \quad (9)$$

then we have, as  $n \rightarrow \infty$ ,

$$\zeta_s(X_n, X) = O(R(n)),$$

where  $\mathcal{L}(X)$  is given as the unique fixed point in  $\mathcal{P}_s^d(0, \text{Id}_d)$  of the equation

$$X \stackrel{d}{=} \sum_{r=1}^K A_r^* X^{(r)} + b^*, \quad (10)$$

with  $(A_1^*, \dots, A_K^*, b^*), X^{(1)}, \dots, X^{(K)}$  independent and  $X^{(r)} \stackrel{d}{=} X$  for  $r = 1, \dots, K$ .

► **Remark 2.** In applications, the convergence rate of the coefficients (conditions (6) and (8)) is often faster than the convergence rate of the quantities  $X_n$ , see, e.g., Section 4.4. In these cases, it is often possible to perform the induction step in the proof of Theorem 1 although condition (9) does not hold. To be more precise, we may assume

$$\|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_s + \|b^{(n)} - b^*\|_s + \|A_r^{(n)} - A_r^*\|_s = O(\tilde{R}(n))$$

for every  $\ell \geq 0$ ,  $r = 1, \dots, K$  and  $n \rightarrow \infty$ . Then, instead of condition (9), it is sufficient to find some  $K > 0$  such that

$$\mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{n_1 \leq I_r^{(n)} < n\}} \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^s \right] \leq 1 - p_n - \frac{\tilde{R}(n)}{KR(n)}$$

for all large  $n$  with  $p_n := \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{I_r^{(n)} = n\}} \|A_r^{(n)}\|_{\text{op}}^s \right]$ .

## 2.2 An improved transfer theorem for normal limit distributions

We now consider the special case where the sequence  $(X_n)_{n \geq 0}$  has finite third moments and satisfies recursion (2) with  $(A_1^{(n)}, \dots, A_K^{(n)}, b^{(n)}) \xrightarrow{L_3} (A_1^*, \dots, A_K^*, b^*)$  for some coefficients  $A_1^*, \dots, A_K^*, b^*$  with finite third moments and

$$b^* = 0, \quad \sum_{r=1}^K A_r^* (A_r^*)^T = \text{Id}_d$$

almost surely. Corollary 3.4 in [14] implies that, if  $\mathbb{E}[\sum_{r=1}^K \|A_r^*\|_{\text{op}}^3] < 1$ , equation (10) has a unique solution in the space  $\mathcal{P}_3^d(0, \text{Id}_d)$ . Furthermore, e.g., using characteristic functions, it is easily checked that this unique solution is the standard normal distribution  $\mathcal{N}(0, \text{Id}_d)$ .

In this special case of normal limit laws, it is possible to derive a refined version of Theorem 1. Instead of the technical condition (6), we now need the weaker condition

$$\|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_3^3 = O(R(n)), \quad n \rightarrow \infty, \quad (11)$$

for all  $\ell \in \mathbb{N}$  and  $r = 1, \dots, K$ . Moreover, condition (8) concerning the convergence rates of the coefficients can be weakened, which is formulated in the following theorem.

► **Theorem 3.** *Let  $(X_n)_{n \geq 0}$  be given as in (5) with finite third moments. We assume that for some  $R(n) > 0$  monotonically decreasing with  $R(n) \rightarrow 0$  as  $n \rightarrow \infty$  we have*

$$\left\| \sum_{r=1}^K A_r^{(n)} (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} + \|b^{(n)}\|_3^3 = O(R(n)), \quad (12)$$

and the technical conditions (7) and (11) being satisfied for  $s = 3$ . If

$$\limsup_{n \rightarrow \infty} \mathbb{E} \sum_{r=1}^K \left( \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^3 \right) < 1, \quad (13)$$

then we have, as  $n \rightarrow \infty$ ,

$$\zeta_3(X_n, \mathcal{N}(0, \text{Id}_d)) = O(R(n)).$$

**Proof.** (Sketch) We define an accompanying sequence  $(Z_n^*)_{n \geq 0}$  by

$$Z_n^* := \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} N^{(r)} + b^{(n)}, \quad n \geq 0,$$

where  $(A_1^{(n)}, \dots, A_K^{(n)}, I^{(n)}, b^{(n)}), N^{(1)}, \dots, N^{(K)}$  are independent,  $\mathcal{L}(N^{(r)}) = \mathcal{N}(0, \text{Id}_d)$  for  $r = 1, \dots, K$  and  $T_n T_n^T = \text{Cov}(X_n)$  for  $n \geq 0$ . Hence,  $Z_n^*$  has a finite third moment,  $\mathbb{E}[Z_n^*] = 0$  and  $\text{Cov}(Z_n^*) = \text{Id}_d$  for all  $n \geq n_1$ . By the triangle inequality, we have

$$\zeta_3(X_n, \mathcal{N}(0, \text{Id}_d)) \leq \zeta_3(X_n, Z_n^*) + \zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)).$$

Then, the assertion follows inductively if one has shown the bound  $\zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)) = O(R(n))$ : Using the convolution property of the multidimensional normal distribution, we obtain the representation

$$Z_n^* = \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} N^{(r)} + b^{(n)} \stackrel{d}{=} G_n N + b^{(n)},$$

where  $G_n G_n^T = \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} T_{I_r^{(n)}}^T (A_r^{(n)})^T$ ,  $\mathcal{L}(N) = \mathcal{N}(0, \text{Id}_d)$  and  $N$  is independent of  $(G_n, b^{(n)})$ . As  $\text{Cov}(Z_n^*) = \text{Id}_d$  for all  $n \geq n_1$ , we have  $\mathbb{E}[G_n G_n^T + b^{(n)}(b^{(n)})^T] = \text{Id}_d$  for  $n \geq n_1$ . Furthermore, we have  $\|b^{(n)}\|_3^3 = O(R(n))$  and

$$\begin{aligned} \|G_n G_n^T - \text{Id}_d\|_{3/2}^{3/2} &= \left\| \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} T_{I_r^{(n)}}^T (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} \\ &= O \left( \left\| \sum_{r=1}^K \mathbf{1}_{\{I_r^{(n)} < n_1\}} A_r^{(n)} (T_{I_r^{(n)}} T_{I_r^{(n)}}^T - \text{Id}_d) (A_r^{(n)})^T \right\|_{3/2}^{3/2} \right. \\ &\quad \left. + \left\| \sum_{r=1}^K A_r^{(n)} (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} \right) \\ &= O \left( \sum_{r=1}^K \|\mathbf{1}_{\{I_r^{(n)} < n_1\}} A_r^{(n)}\|_3^3 + \left\| \sum_{r=1}^K A_r^{(n)} (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} \right) \\ &= O(R(n)). \end{aligned}$$

Thus, the following Lemma 5 implies  $\zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)) = O(R(n))$ . Lemma 5 is the main part of the present proof.  $\blacktriangleleft$

► **Remark 4.** Theorem 3, when applicable, often improves over Theorem 1 by a factor 3 in the exponent, see Remark 9 for an example. This is caused by the additional exponents in (12) in comparison to (8).

► **Lemma 5.** *Let  $(Z_n^*)_{n \geq 0}$  be a sequence of  $d$ -dimensional random vectors satisfying  $Z_n^* \stackrel{d}{=} G_n N + b^{(n)}$  with some random  $(d \times d)$ -matrix  $G_n$  and some random vector  $b^{(n)}$  such that  $\mathbb{E}[Z_n^*] = 0$ ,  $\text{Cov}(Z_n^*) = \text{Id}_d$  and  $N \sim \mathcal{N}(0, \text{Id}_d)$  is independent of  $(G_n, b^{(n)})$ . Furthermore, we assume that, as  $n \rightarrow \infty$ ,*

$$\|G_n G_n^T - \text{Id}_d\|_{3/2}^{3/2} + \|b^{(n)}\|_3^3 = O(R(n))$$

for appropriate  $R(n)$ . Then, we have, as  $n \rightarrow \infty$ ,

$$\zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)) = O(R(n)).$$

The proof of Lemma 5 builds upon ideas of [15].

### 3 Expansions of moments

In applications to problems arising in theoretical computer science, where the recurrence (1) is explicitly given, one usually has no direct means to identify the orders of the terms  $\|b^{(n)} - b^*\|_s$  and  $\|A_r^{(n)} - A_r^*\|_s$ . This is due to the fact that the mean vector  $M_n$  and the covariance matrix  $C_n$ , for the cases  $1 < s \leq 2$  and  $2 < s \leq 3$  respectively, which are used for the normalization (5) are typically not exactly known or too involved to be amenable to explicit calculations. As a substitute one usually has asymptotic expansions of these sequences as  $n \rightarrow \infty$ .

In the present section we assume the dimension to be  $d = 1$  and  $A_r(n) = 1$  for all  $r = 1, \dots, K$  and provide tools to apply the general Theorems 1 and 3 on the basis of expansions of the mean and variance. We assume that

$$\mathbb{E}[X_n] = \mu(n) = f(n) + O(e(n)), \quad \text{Var}(X_n) = \sigma^2(n) = g(n) + O(h(n)), \quad (14)$$

with  $e(n) = o(f(n))$  and  $h(n) = o(g(n))$ . To connect Theorems 1 and 3 to recurrences with known expansions we use the following notion.

► **Definition 6.** A sequence  $(a(n))_{n \geq 0}$  of non-negative numbers is called essentially non-decreasing if there exists a  $c > 0$  such that  $a(m) \leq ca(n)$  for all  $0 \leq m < n$ .

The scaling introduced in (5) with the special choices  $A_r(n) = 1$  for all  $r = 1, \dots, K$  leads to the scaled recurrence for  $(X_n)$  given in (2) with

$$A_r^{(n)} = \frac{\sigma(I_r^{(n)})}{\sigma(n)}, \quad b^{(n)} = \frac{1}{\sigma(n)} \left( b_n - \mu(n) + \sum_{r=1}^K \mu(I_r^{(n)}) \right). \quad (15)$$

Additionally, we consider the corresponding quantities

$$\bar{A}_r^{(n)} = \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)}, \quad \bar{b}^{(n)} = \frac{1}{g^{1/2}(n)} \left( b_n - f(n) + \sum_{r=1}^K f(I_r^{(n)}) \right). \quad (16)$$

Then we have:

► **Lemma 7.** With  $A_r^{(n)}$ ,  $b^{(n)}$  given in (15),  $\bar{A}_r^{(n)}$ ,  $\bar{b}^{(n)}$  given in (16), and the expansions for  $\mu(n)$ ,  $\sigma^2(n)$  given in (14) the following holds.

If the sequence  $h/g^{1/2}$  is essentially non-decreasing then

$$\|A_r^{(n)} - A_r^*\|_s \leq \|\bar{A}_r^{(n)} - A_r^*\|_s + O\left(\frac{h(n)}{g(n)}\right). \quad (17)$$

If the sequence  $h$  is essentially non-decreasing then

$$\left\| \sum_{r=1}^K (A_r^{(n)})^2 - 1 \right\|_s \leq \left\| \sum_{r=1}^K (\bar{A}_r^{(n)})^2 - 1 \right\|_s + O\left(\frac{h(n)}{g(n)}\right). \quad (18)$$

If the sequence  $e$  is essentially non-decreasing then

$$\|b^{(n)} - b^*\|_s \leq \|\bar{b}^{(n)} - b^*\|_s + O\left(\frac{h(n)}{g(n)} + \frac{e(n)}{g^{1/2}(n)}\right). \quad (19)$$

If the sequence  $g/h$  is essentially non-decreasing and

$$T(n) := \mathbb{E} \sum_{r=1}^K \frac{g^{s/2-1}(I_r^{(n)})h(I_r^{(n)})R(I_r^{(n)})}{g^{s/2}(n)R(n)}$$

then we have

$$\mathbb{E} \sum_{r=1}^K \frac{\sigma^s(I_r^{(n)})R(I_r^{(n)})}{\sigma^s(n)R(n)} \leq \mathbb{E} \sum_{r=1}^K \frac{g^{s/2}(I_r^{(n)})R(I_r^{(n)})}{g^{s/2}(n)R(n)} + O(T(n)). \quad (20)$$

**Proof.** We show (17), the other bounds can be shown similarly. Note that  $\sigma^2(n) = g(n) + O(h(n))$  implies  $\sigma(n) = g^{1/2}(n) + O(h(n)/g^{1/2}(n))$  and that for any essentially non-decreasing sequence  $(a(n))_{n \geq 0}$  we have  $\|a(I_r^{(n)})\|_\infty = O(a(n))$ . Since  $h/g^{1/2}$  is essentially non-decreasing we obtain

$$\begin{aligned} A_r^{(n)} &= \frac{\sigma(I_r^{(n)})}{\sigma(n)} = \frac{g^{1/2}(I_r^{(n)}) + O(h(I_r^{(n)})/g^{1/2}(I_r^{(n)}))}{\sigma(n)} \\ &= \frac{g^{1/2}(I_r^{(n)}) + O(h(n)/g^{1/2}(n))}{g^{1/2}(n)} \cdot \frac{g^{1/2}(n)}{\sigma(n)} \\ &= \left( \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)} + O\left(\frac{h(n)}{g(n)}\right) \right) \left( 1 + O\left(\frac{h(n)}{g(n)}\right) \right) \\ &= \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)} + O\left(\frac{h(n)}{g(n)} \left( 1 + \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)} \right)\right). \end{aligned}$$

Hence, we obtain

$$\|A_r^{(n)} - A_r^*\|_s \leq \|\bar{A}_r^{(n)} - A_r^*\|_s + O\left(\frac{h(n)}{g(n)} \left( 1 + \|\bar{A}_r^{(n)}\|_s \right)\right).$$

Since  $\bar{A}_r^{(n)} \rightarrow A_r^*$  in  $L_s$  we have  $\|\bar{A}_r^{(n)}\|_s = O(1)$ , hence

$$\|A_r^{(n)} - A_r^*\|_s \leq \|\bar{A}_r^{(n)} - A_r^*\|_s + O\left(\frac{h(n)}{g(n)}\right),$$

which is bound (17). ◀

Note that in applications the terms on the right hand side in the estimates (17)–(20) can easily be bounded when expansions as in (14) with explicit functions  $e, f, g, h$  are available.

## 4 Applications

We start by deriving a known result to illustrate in detail how to apply our framework of the previous sections.

### 4.1 Quicksort: Key comparisons

The number of key comparisons  $Y_n$  needed by the Quicksort algorithm to sort  $n$  randomly permuted (distinct) numbers satisfies the distributional recursion

$$Y_n \stackrel{d}{=} Y_{I_n} + Y'_{n-1-I_n} + n - 1, \quad n \geq 1, \tag{21}$$

where  $Y_0 := 0$  and  $(Y_k)_{k=0, \dots, n-1}, (Y'_k)_{k=0, \dots, n-1}, I_n$  are independent,  $I_n$  is uniformly distributed on  $\{0, \dots, n-1\}$ , and  $Y_k \stackrel{d}{=} Y'_k, k \geq 0$ . Hence, equation (21) is covered by our general recurrence (1). For the expectation and variance of  $Y_n$  exact expressions are known which imply the asymptotic expansions

$$\begin{aligned} \mathbb{E}Y_n &= 2n \log(n) + (2\gamma - 4)n + O(\log n), \\ \text{Var}(Y_n) &= \sigma^2 n^2 - 2n \log(n) + O(n), \end{aligned}$$

where  $\gamma$  denotes Euler's constant and  $\sigma := \sqrt{7 - 2\pi^2/3} > 0$ . We introduce the normalized quantities  $X_0 := X_1 := X_2 := 0$  and

$$X_n := \frac{Y_n - \mathbb{E}Y_n}{\sqrt{\text{Var}(Y_n)}}, \quad n \geq 3.$$

To apply Theorem 1 we need to find an  $0 < s \leq 3$  and a sequence  $(R(n))$  with (8) and (9). Note that the  $Y_n$  are bounded, thus  $L_s$ -integrable for any  $s > 0$ . To bound the  $L_s$ -norms appearing in (8) we use Lemma 7 and choose

$$\begin{aligned} f(n) &= 2n \log(n) + (2\gamma - 4)n, & e(n) &= \log n, \\ g(n) &= \sigma^2 n^2, & h(n) &= n \log n. \end{aligned}$$

With these functions we obtain for the quantities defined in (16) that

$$\begin{aligned} \bar{A}_1^{(n)} &= \frac{I_n}{n}, & \bar{A}_2^{(n)} &= \frac{n-1-I_n}{n}, \\ \bar{b}^{(n)} &= \frac{1}{\sigma} \left( 2 \frac{I_n}{n} \log \frac{I_n}{n} + 2 \frac{n-1-I_n}{n} \log \frac{n-1-I_n}{n} + \frac{n-1}{n} + O\left(\frac{\log n}{n}\right) \right). \end{aligned}$$

With the embedding  $I_n = \lfloor nU \rfloor$  with  $U$  uniformly distributed over the unit interval  $[0, 1]$  we have

$$A_1^* = U, \quad A_2^* = 1 - U, \quad b^* = \frac{1}{\sigma} (2U \log(U) + 2(1-U) \log(1-U) + 1) =: \frac{1}{\sigma} \varphi(U).$$

The limit theorem  $X_n \rightarrow X$  has been derived by different methods by Régnier [16] and Rösler [17]. Rösler [17] also found that the scaled limit  $Y := \sigma X$  satisfies the distributional fixed-point equation

$$Y \stackrel{d}{=} UY + (1-U)Y' + \varphi(U).$$

Lower and upper bounds for the rate of convergence in  $X_n \rightarrow X$  have been studied for various metrics in Fill and Janson [6] and Neininger and Rüschemdorf [13].

Now, we apply the framework of the present paper: For  $r = 1, 2$  and any  $s \geq 1$  we find that

$$\|\bar{A}_r^{(n)} - A_r^*\|_s = O\left(\frac{1}{n}\right).$$

Using Proposition 3.2 of Rösler [17] we obtain

$$\|\bar{b}^{(n)} - b^*\|_s = O\left(\frac{\log n}{n}\right).$$

Moreover, we have

$$\frac{h(n)}{g(n)} = O(R(n)) \quad \text{and} \quad \frac{e(n)}{g^{1/2}(n)} = O(R(n)) \quad \text{with} \quad R(n) := \frac{\log n}{n},$$

thus Lemma 7 implies that condition (8) is satisfied for our choice of the sequence  $R$ . To verify condition (9) by use of (20) we obtain that for  $T(n)$  given in Lemma 7 we find  $T(n) = O(\log(n)/n) \rightarrow 0$  and that

$$\mathbb{E} \sum_{r=1}^2 \frac{g^{s/2}(I_r^{(n)}) R(I_r^{(n)})}{g^{s/2}(n) R(n)} = \mathbb{E} \sum_{r=1}^2 \left( \frac{I_r^{(n)}}{n} \right)^{s-1} \frac{\log I_r^{(n)}}{\log n}.$$

Note that the latter expression has a limit superior of less than 1 if and only if  $s > 2$ . Hence, Theorem 1 is applicable for  $s > 2$  and yields that

$$\zeta_s(X_n, X) = O\left(\frac{\log n}{n}\right), \quad \text{for } 2 < s \leq 3. \quad (22)$$

The bound (22) had previously been shown for  $s = 3$  in [13], where also the optimality of the order was shown, i.e., that  $\zeta_3(X_n, X) = \Theta(\log(n)/n)$ .

In the planned full paper version we also discuss bounds on rates of convergence for various cost measures of the related Quickselect algorithms under various models for the rank to be selected.



### 4.2 Size of $m$ -ary search trees

The size of  $m$ -ary search trees satisfies the recurrence (1) with  $K = m \geq 3$ ,  $A_1(n) = \dots = A_m(n) = 1$ ,  $n_0 = m$ ,  $b_n = 1$ , i.e., we have

$$Y_n \stackrel{d}{=} \sum_{r=1}^m Y_{I_r^{(n)}}^{(r)} + 1, \quad n \geq m.$$

For a representation of  $I^{(n)}$  we define for independent, identically  $\text{unif}[0, 1]$  distributed random variables  $U_1, \dots, U_{m-1}$  their spacings in  $[0, 1]$  by  $S_1 = U_{(1)}, S_2 = U_{(2)} - U_{(1)}, \dots, S_m := 1 - U_{(m-1)}$ , where  $U_{(1)}, \dots, U_{(m-1)}$  denote the order statistics of  $U_1, \dots, U_{m-1}$ . Then  $I^{(n)}$  has the mixed multinomial distribution:

$$I^{(n)} \stackrel{d}{=} M(n - m + 1, S_1, \dots, S_m).$$

By this we mean that given  $(S_1, \dots, S_m) = (s_1, \dots, s_m)$  we have that  $I^{(n)}$  is multinomial  $M(n - m + 1, s_1, \dots, s_m)$  distributed. Expectations, variances and limit laws for  $Y_n$  have been studied, see [12, 4]. We have

$$\mathbb{E}Y_n = \mu n + O(1 + n^{\alpha-1}), \quad m \geq 3, \tag{23}$$

$$\text{Var}(Y_n) = \sigma^2 n + O(1 + n^{2\alpha-2}), \quad 3 \leq m \leq 26, \tag{24}$$

Here, the constants  $\mu, \sigma > 0$  depend on  $m$  and  $\alpha \in \mathbb{R}$  depends on  $m$  such that  $\alpha < 1$  for  $m \leq 13$ ,  $1 \leq \alpha \leq 4/3$  for  $14 \leq m \leq 19$ , and  $4/3 \leq \alpha \leq 3/2$  for  $20 \leq m \leq 26$ , see, e.g., Mahmoud [12, Table 3.1] for the values  $\alpha = \alpha_m$  depending on  $m$ . It is known that  $Y_n$  standardized by mean and variance satisfies a central limit law for  $m \leq 26$ , whereas the standardized sequence has no weak limit for  $m > 26$  due to dominant periodicities, see Chern and Hwang [4]. The rate of convergence in the central limit law for  $m \leq 26$  for the Kolmogorov metric has been identified in Hwang [9]. Our Theorem 3 implies the central limit theorem for  $Y_n$  with  $m \leq 26$  with the same (up to an  $\varepsilon$  for  $3 \leq m \leq 19$ ) rate of convergence for the Zolotarev metric  $\zeta_3$ :

► **Theorem 8.** *The size  $Y_n$  of a random  $m$ -ary search tree with  $n$  items inserted satisfies, for  $m \leq 26$  and any  $\varepsilon > 0$ ,*

$$\zeta_3\left(\frac{Y_n - \mathbb{E}Y_n}{\sqrt{\text{Var}(Y_n)}}, \mathcal{N}(0, 1)\right) = \begin{cases} O(n^{-1/2+\varepsilon}), & 3 \leq m \leq 19, \\ O(n^{-3(3/2-\alpha)}), & 20 \leq m \leq 26, \end{cases}$$

as  $n \rightarrow \infty$ .

**Proof.** In order to apply Theorem 3 we have to estimate the orders of  $\|\sum_{r=1}^m (A_r^{(n)})^2 - 1\|_{3/2}$  and  $\|b^{(n)}\|_3$  with  $A_r^{(n)}$  and  $b^{(n)}$  defined in (3). For this we apply Lemma 7. From (23) and (24) we obtain that for the quantities appearing in Lemma 7 we can choose  $f(n) = \mu n$ ,  $e(n) = 1 \vee n^{\alpha-1}$ ,  $g(n) = \sigma^2 n$ , and  $h(n) = 1 \vee n^{2(\alpha-1)}$ . Hence we obtain

$$\left\| \sum_{r=1}^m (\overline{A}_r^{(n)})^2 - 1 \right\|_{3/2} = \left\| \sum_{r=1}^m \frac{I_r^{(n)}}{n} - 1 \right\|_{3/2} = \frac{m-1}{n} = O(n^{-1})$$

and  $O(h(n)/g(n)) = O(n^{-(1 \wedge (3-2\alpha))})$ . This implies

$$\left\| \sum_{r=1}^m (A_r^{(n)})^2 - 1 \right\|_{3/2}^{3/2} = O(n^{-((3/2) \wedge (3(3/2-\alpha)))}).$$

Similarly we obtain

$$\|\bar{b}^{(n)}\|_3 = \frac{1}{\sigma\sqrt{n}} \left\| 1 - \mu n + \sum_{r=1}^m \mu I_r^{(n)} \right\|_3 = \frac{1}{\sigma\sqrt{n}} \|1 - \mu(m-1)\|_3 = O(n^{-1/2})$$

and  $O(e(n)/g^{1/2}(n)) = O(n^{-(1/2 \wedge (3/2 - \alpha))})$ . This implies

$$\|b^{(n)}\|_3^3 = O(n^{-((3/2) \wedge (3(3/2 - \alpha)))}).$$

Hence, condition (12) is satisfied with  $R(n) = n^{-((3/2) \wedge (3(3/2 - \alpha)))}$ . ◀

► **Remark 9.** Using Theorem 1 instead of Theorem 3 in the latter proof is also possible but leads to a bound  $O(n^{-(3/2 - \alpha)})$  for  $20 \leq m \leq 26$ , missing the factor 3 appearing in Theorem 8.

In the full paper version we also discuss rates of convergence for the number of leaves of  $d$ -dimensional random point quadtrees in the model of [7, 3, 8] where a similar behavior as in Theorem 8 appears. A technically related example is the number of maxima in right triangles in the model of [1, 2], where the order  $n^{-1/4}$  appears. Our framework also applies.

### 4.3 Periodic functions in mean and variance

We now discuss some applications where the asymptotic expansions of the mean and the variance include periodic functions instead of fixed constants. This is the case for several quantities in binomial splitting processes such as tries, PATRICIA tries and digital search trees. Throughout this section, we assume that we have a sequence  $(Y_n)_{n \geq 0}$  with finite third moments satisfying the recursion

$$Y_n \stackrel{d}{=} Y_{I_1^{(n)}}^{(1)} + Y_{I_2^{(n)}}^{(2)} + b_n, \quad n \geq n_0, \tag{25}$$

with  $(I^{(n)}, b_n)$ ,  $(Y_n^{(1)})_{n \geq 0}$  and  $(Y_n^{(2)})_{n \geq 0}$  independent and  $(Y_n^{(r)})_{n \geq 0} \stackrel{d}{=} (Y_n)_{n \geq 0}$  for  $r = 1, 2$ . Furthermore,  $I_1^{(n)}$  has the binomial distribution  $\text{Bin}(n, \frac{1}{2})$  and  $I_2^{(n)} = n - I_1^{(n)}$  or  $I_1^{(n)}$  is binomially  $\text{Bin}(n-1, \frac{1}{2})$  distributed and  $I_2^{(n)} = n-1 - I_1^{(n)}$ . Mostly, these binomial recurrences are asymptotically normally distributed, see [10, 11, 14, 18] for some examples.

Our first theorem covers the case of linear mean and variance, i.e. we assume that, as  $n \rightarrow \infty$ ,

$$\mathbb{E}[Y_n] = nP_1(\log_2 n) + O(1), \tag{26}$$

$$\text{Var}(Y_n) = nP_2(\log_2 n) + O(1), \tag{27}$$

for some smooth and 1-periodic functions  $P_1, P_2$  with  $P_2 > 0$ . Possible applications would start with the analysis of the number of internal nodes of a trie for  $n$  strings in the symmetric Bernoulli model and the number of leaves in a random digital search tree, see, e.g., [10].

► **Theorem 10.** *Let  $(Y_n)_{n \geq 0}$  have finite third moments and satisfy (25) with  $\|b_n\|_3 = O(1)$ , (26) and (27). Then, for any  $\varepsilon > 0$  and  $n \rightarrow \infty$ , we have*

$$\zeta_3\left(\frac{Y_n - \mathbb{E}[Y_n]}{\sqrt{\text{Var}(Y_n)}}, \mathcal{N}(0, 1)\right) = O(n^{-1/2 + \varepsilon}).$$

We now consider the case where our quantities  $Y_n$  satisfy recursion (25) with  $b_n$  being essentially  $n$ . We assume that, as  $n \rightarrow \infty$ , we have

$$\mathbb{E}[Y_n] = n \log_2(n) + nP_1(\log_2 n) + O(1), \tag{28}$$

$$\text{Var}(Y_n) = nP_2(\log_2 n) + O(1), \tag{29}$$

for some smooth and 1-periodic functions  $P_1, P_2$  with  $P_2 > 0$ . This covers, for example, the external path length of random tries and related digital tree structures constructed from  $n$  random binary strings under appropriate independence assumptions.

► **Theorem 11.** *Let  $(Y_n)_{n \geq 0}$  have finite third moments and satisfy (25) with  $\|b_n - n\|_3 = O(1)$ , (28) and (29). Then, for any  $\varepsilon > 0$  and  $n \rightarrow \infty$ , we have*

$$\zeta_3\left(\frac{Y_n - \mathbb{E}[Y_n]}{\sqrt{\text{Var}(Y_n)}}, \mathcal{N}(0, 1)\right) = O(n^{-1/2+\varepsilon}).$$

### 4.4 A multivariate application

We consider a random binary search tree with  $n$  nodes built from a random permutation of  $\{1, \dots, n\}$ . For  $n \geq 0$ , we denote by  $L_{0n}$  the number of nodes with no left descendant and by  $L_{1n}$  the number of nodes with exactly one left descendant. Defining  $Y_n := (L_{0n}, L_{1n})$ , we have  $Y_0 = (0, 0)$  and we obtain the following distributional recurrence:

$$Y_n \stackrel{d}{=} Y_{I_1^{(n)}}^{(1)} + Y_{I_2^{(n)}}^{(2)} + b_n, \quad n \geq 1,$$

where  $(Y_j^{(1)})_{j \geq 0}$  and  $(Y_j^{(2)})_{j \geq 0}$  are independent copies of  $(Y_j)_{j \geq 0}$ ,  $I_1^{(n)}$  is uniformly distributed on  $\{0, \dots, n-1\}$  and independent of  $(Y^{(1)})$  and  $(Y^{(2)})$ ,  $I_2^{(n)} = n-1-I_1^{(n)}$  and  $b_n = (\mathbf{1}_{\{I_1^{(n)}=0\}}, \mathbf{1}_{\{I_1^{(n)}=1\}})$ . In Devroye [5] it is shown that, for  $n \geq 2$ ,

$$\mathbb{E}[L_{0n}] = \frac{1}{2}(n+1), \quad \mathbb{E}[L_{1n}] = \frac{1}{6}(n+1),$$

and that the standardized quantities have a limiting normal distribution. Using Devroye’s description with local counters one also obtains the covariance structure:

► **Lemma 12.** *For  $n \geq 4$ , we have  $\text{Cov}(Y_n) = (n+1)\Gamma$  with*

$$\Gamma = \frac{1}{360} \begin{pmatrix} 30 & -15 \\ -15 & 28 \end{pmatrix}.$$

For  $n \geq 0$ , we now set  $M_n := \mathbb{E}[Y_n]$ ,  $C_n = \text{Id}_2$  for  $n \leq 3$ ,  $C_n := \text{Cov}(Y_n)$  for  $n \geq 4$  and define  $X_n := C_n^{-1/2}(Y_n - M_n)$  for  $n \geq 0$ . Note that the matrix  $\Gamma$  in Lemma 12 is symmetric and positive definite, which implies, for  $n \geq 4$ ,

$$C_n^{1/2} = \sqrt{n+1} \Gamma^{1/2} \quad \text{and} \quad C_n^{-1/2} = \frac{1}{\sqrt{n+1}} \Gamma^{-1/2}.$$

The normalized quantities satisfy  $X_0 = (0, 0)$  and recursion (2) with  $K = 2$ ,  $n_0 = 1$ ,

$$A_r^{(n)} = C_n^{-1/2} C_{I_r^{(n)}}^{1/2} = \mathbf{1}_{\{I_r^{(n)} \geq 4\}} \sqrt{\frac{I_r^{(n)} + 1}{n+1}} \text{Id}_2 + \mathbf{1}_{\{I_r^{(n)} < 4\}} \frac{1}{\sqrt{n+1}} \Gamma^{-1/2}$$

for  $r = 1, 2$  and

$$b^{(n)} = C_n^{-1/2}(b_n - M_n + M_{I_1^{(n)}} + M_{I_2^{(n)}}).$$

Modeling all quantities on a joint probability space such that  $I_1^{(n)}/n$  converges almost surely to a uniform random variable  $U$  in  $[0, 1]$ , we have the  $L_3$ -convergences  $A_1^{(n)} \rightarrow \sqrt{U} \text{Id}_2$ ,  $A_2^{(n)} \rightarrow \sqrt{1-U} \text{Id}_2$  and  $b^{(n)} \rightarrow 0$  as  $n \rightarrow \infty$ . Thus, we are in the situation of Section 2.2 and obtain the limiting equation

$$X \stackrel{d}{=} \sqrt{U}X^{(1)} + \sqrt{1-U}X^{(2)},$$

with  $U$  uniformly distributed on  $[0, 1]$  and  $X^{(1)}, X^{(2)}$  and  $U$  independent. We now check the conditions of Theorem 3. Since  $A_1^{(n)}(A_1^{(n)})^T + A_2^{(n)}(A_2^{(n)})^T = \text{Id}_2$  on the event  $\{I_1^{(n)}, I_2^{(n)} \geq 4\}$ , we obtain, as  $n \rightarrow \infty$ ,

$$\begin{aligned} \left\| \sum_{r=1}^2 A_r^{(n)}(A_r^{(n)})^T - \text{Id}_2 \right\|_{3/2}^{3/2} &= O\left( \left\| \mathbf{1}_{\{I_1^{(n)} < 4\}} \left( \frac{1}{n+1} \Gamma^{-1} + \frac{I_2^{(n)} + 1}{n+1} \text{Id}_2 - \text{Id}_2 \right) \right\|_{3/2}^{3/2} \right) \\ &= O\left( \mathbb{E} \left[ \mathbf{1}_{\{I_1^{(n)} < 4\}} \left\| \frac{1}{n+1} \Gamma^{-1} - \frac{I_1^{(n)} + 1}{n+1} \text{Id}_2 \right\|_{\text{op}}^{3/2} \right] \right) \\ &= O(n^{-5/2}). \end{aligned}$$

Similarly, we obtain

$$\|b^{(n)}\|_3^3 = O(n^{-5/2}).$$

Since we have  $\|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_3^3 = O(n^{-5/2})$  for  $\ell \in \mathbb{N}$  and  $r = 1, 2$ , the technical conditions are satisfied. We now use Theorem 3 with  $R(n) = n^{-1/2}$ . Note that condition (13) is not satisfied for  $R(n) = n^{-1/2}$ , but we can use the weakened condition stated in Remark 2 to obtain the following result.

► **Theorem 13.** Denoting by  $Y_n := (L_{0n}, L_{1n})$  the vector of the numbers of nodes with no and with exactly one left descendant respectively in a random binary search tree with  $n$  nodes we have, for  $n \rightarrow \infty$ , that

$$\zeta_3(\text{Cov}(Y_n)^{-1/2}(Y_n - \mathbb{E}[Y_n]), \mathcal{N}(0, \text{Id}_2)) = O(n^{-1/2}).$$

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